## Double DES and Triple DES For IT $7^{\text {th }}$ Sem Students

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## Double DES



## Intermediate CT

- Mathematically
$\mathrm{Ic}=\mathrm{Ek} 1(\mathrm{P})$
k1
$\mathrm{CT}=\operatorname{Ek} 2(\operatorname{Ek} 1(\mathrm{P}))$


## k2 <br> Encryption Process

## Cont...



Intermediate CT


Decryption Process

## Problem of Double DES

- Markel and Hellman introduced encryption from one end and decryption from other end and matching the results in the middle hence the name "meet in the middle attack".


## Meet in the Middle Attack

- Suppose that cryptanalysis knows two basic pieces of information P (a plain text block) and CT(corresponding the final cipher text block) for a message.


## Temporary



## Cont...

- The result of $1^{\text {st }}$ encryption is called as T and denoted

$$
\mathrm{T}=\mathrm{Ek} 1(\mathrm{P})
$$

- After this encryption the encrypted block is encrypted with another key k2 then

$$
\mathrm{CT}=\mathrm{Ek} 2(\operatorname{Ek} 1(\mathrm{P}))
$$

- Now the aim of the cryptanalysis who is armed with the knowledge of P and C is to obtain the values of k 1 and k 2 the cryptanalysis do...


## Cont...

- Step 1- for all possible values of $2^{\wedge} 56$ of $k 1$ the cryptanalysis would use a large table in the memory of the computer and perform the following two points...
- 1-the cryptanalysis would encrypt the plain text block P by performing the $1^{\text {st }}$ encryption operation.

$$
\text { i.e. } T=\operatorname{Ek} 1(\mathrm{P})
$$

- 2-the cryptanalysis store the output of the operation Ek1(P) in temporary T and calculate

$$
\mathrm{CT}=\operatorname{Ek} 2(\operatorname{Ek} 1(\mathrm{P}))
$$

## Cont...

- Step 2- for decryption process

$$
\begin{gathered}
\mathrm{T}=\mathrm{Dk} 2(\mathrm{CT}) \\
\mathrm{PT}=\operatorname{Dk} 1(\mathrm{Dk} 2(\mathrm{P}))
\end{gathered}
$$

- From above two steps

$$
\mathrm{T}=\operatorname{Ek} 1(\mathrm{P})=\mathrm{Dk} 2(\mathrm{CT})
$$

- Now if the cryptanalysis creates a table of $\operatorname{Ek} 1(\mathrm{P})$ for all possible values of k 1 and then perform Dk2(CT) for all possible values of $k 2$,so there is a chance that she or he gets the same T in both operation.


## Cont...

- If the cryptanalysis is able to find the same T for both encryption with k 1 and decryption with k2,its means that the cryptanalysis knows not only P and C but he has been also able to find out the possible values of k 1 and k 2 .


## Triple DES

- Although the meet in the middle attack on double DES is not quite practical yet in cryptography, but it is always better to minimum chances.
- As we can imagine triple DES is DES three times. It comes in two variations like...
- Triple DES with Three keys.
- Triple DES with two keys.


## Triple DES with Three Keys



CT=Ek3(Ek2(Ek1(PT)))

## Triple DES with Two Keys



## Reference

- Cryptography and network security "Atul Kahate" 3e,Mc Graw hill education.

